A network traffic sources modeling method based on measured data defragmentation

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Abstract- Over recent years, the need for simulating complex communication networks, in order to assist evaluation, construction, and the upgrading of communication networks has become a key element in optimizing the exploitation of particular networks. In this regard traffic modeling has a key impact on network simulation reliability and, consequently, usability. For these reasons, we have developed and compared two algorithms of traffic source modeling. Both algorithms are based on mimicking a defragmentation process. They come from the captured packet estimated parameters of the probability density function regarding data files sizes and files interarrival times processes. The first one uses in-depth analysis of packet headers for estimating of data file lengths, and the second one does this by measuring packets' lengths and identifying the source and destination IP addresses. Both algorithms consider a TCP/IP encapsulation process.

Keywords- network traffic, traffic modeling, traffic simulation, statistic parameters' estimation.

I. INTRODUCTION

Self similar network traffic [1, 2, 4, 5] is usually modeled from an application point of view. It is usually supposed that the statistics of file sizes and file inter-arrival times are a known [3]. Such kinds of traffic models are supported by most commercial telecommunication simulation tools such as the OPNET Modeler [14, 16, 17, 19, 20, 21], as used in our simulations and experiments. Since packets are generated from files by an encapsulation process, in those cases where files are larger than a packet payload, they are fragmented so that packets are no bigger than the Maximal Transmission Unit (MTU) size. Consequently, for using the measured data of packet traffic when modeling file statistics, it is necessary to transform packet statistics into file statistics [9, 10]. This transformation contains opposite operations to the fragmentation and encapsulation process. The file distribution parameters can be estimated from histograms of transformed statistics by probability density function (pdf) fitting tools [11, 12, 13, 22] or methods, such as CCDF [5, 6] or Hill's estimator [15].

II. PACKET TRAFIC STATISTIC

The self-similar network packet traffic [5, 6, 18] $Z_p(t)$, which usually poses long-range dependence [7, 8], can

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be generally described using a combination of two stochastic processes:

$$Z_p(t) = \psi \left\{ X_p(t), Y_p(t) \right\}$$
(1)

where ψ is the function which depends on the packetsize process $X_p(t)$ and inter-arrival time process $Y_p(t)$. These processes are generated in a higher layer of the TCP/IP reference model from equivalent processes $X_f(t)$ and $Y_f(t)$ performed in data sources (Fig. 1). These layers at data transmission, perform fragmentation and encapsulation of data files into packets and, on reception, the defragmentation and decapsulation of packets into data files:

$$X_{m}(t) \xleftarrow{fragmentation \\ encapsulation}}_{decapsulation} X_{f}(t)$$
(2)

$$Y_m(t) \xleftarrow{fragmentation \\ encapsulation}{\underbrace{decapsulation \\ defragmentation}} Y_f(t)$$
(3)

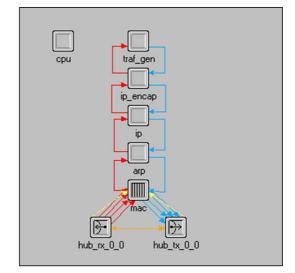


Figure 1: Node model for used IP station in simulation

Let us suppose, that the measured $Z_m(t)$, and modeled $Z_s(t)$ packet traffics are statistically equal, i.e.:

$$Z_m(t) \approx Z_s(t) , \qquad (4)$$

where symbol \approx is used for statistical equivalence, then also constituent of the Z process, i.e. X and Y of measured and modeled packet traffic had to be statistically equal:

$$X_m(t) \approx X_s(t)$$
 and $Y_m(t) \approx Y_s(t)$ (6)

Since $Z_s(t) = \psi \{ X_s(t), Y_s(t) \}$ is generated from $Z_f(t)$ by a fragmentation process, then the defragmentation of $Z_m(t)$ gives $Z_f(t)$, which is the searched for result.

There are many possibilities how estimating files' lengths and their inter-arrival times at the application layer of the communication model. We investigated and compared the results of two algorithms:

- 1. algorithm with in-depth analysis of all packet headers,
- 2. algorithm with coarse inspection of IP header only.

III. PACKETS' STATISTICS TRANSFORMATION

Both algorithms calculate histograms of file source statistics. The main differences between them are complexity and the needed execution time. The first algorithm mimics a complete decapsulation process, and defragmentation in higher layers of the communication model; the second skips decapsulation by considering the average lengths of packet headers and then uses only packet lengths and inter-arrival times. These differences cause differently obtained results, but the differences are negligible, as analyzed later.

A. First algorithm

Each packed IP header has four so called fragmentation fields containing information about data fragmentation (Fig. 2).

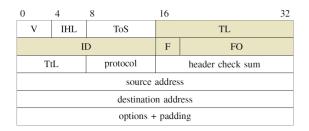


Figure 2: IP header. Shadowed fields are used in the defragmentation process. *Legend*: V: protocol version; IHL: Internet Header Length; ToS: Type of Service; TL: Total Length; ID: Identification Data; F: Flags; FO: Fragment Offset; TtL: Time to Live.

Any sniffers are able to extract these data from the IP header. Knowing them, it is then simple to calculate a length of IP PDU (Protocol Data Unit) which also contains a header of higher layer protocols. Using in-depth header analysis, it is possible, in the similar way to the IP header, calculate the lengths of all these headers.

The first algorithm (Fig. 4) completely respects RFC 793 [26]. It calculates exactly the lengths of the source files, and sorts these lengths within a histogram of the process $X_f(t)$.

From the time stampings of the first and last packets originated from the same file, it calculates the interarrival times and sorts the results within the histogram of the process $Y_f(t)$.

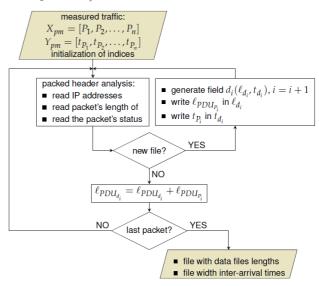


Figure 4: Simplified flowchart of the first algorithm [24].

B. Second algorithm

The second algorithm (Fig. 5), from all the data contained in the packets' headers, uses only source and destination IP addresses. In addition, it uses the time stamping and packets' length data provided by the sniffer. The step from histograms of $X_f(t)$ and $Y_f(t)$ to the parameters of the pdf of these processes is made by EasyFit tool [26].

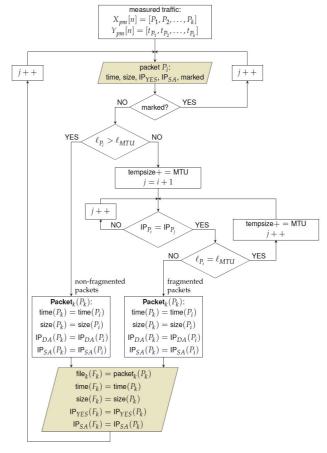


Figure 5: Flowchart of the second algorithm [25].

The second algorithm work as follows. If the packet with the new IP source and/or destination address is the result of fragmentation, i.e. $\ell_{PDU_i} = \ell_{MTU}$, then an algorithm investigates all packets with the same destination and source IP address, in sequence, to those first packets shorter than ℓ_{MTU} . This algorithm considers packets that belong to the same file and sums up their lengths and the length of the originating file, and from them subtracts the average lengths of the IP and TCP headers of each packet (Fig. 6).

		Capit	ired traff	$C Z_m(t)$		
No.	Time	Packet s	size So	ource	Destination	
1	0.0000	70	192.1	68.1.1 1	192.168.1.2	
2	0.1000	1502	192.1	68.1.1 1	192.168.1.2	
3	0.1030	1502	192.1	68.1.1 1	192.168.1.2	
4	0.1050	1502	192.1	68.1.1 1	192.168.1.2	
5	0.1090	540	192.1	68.1.1 1	192.168.1.2	
6	0.3000	65	192.1	68.1.1 1	192.168.1.2	
7	0.5000	78	192.1	68.1.1 1	192.168.1.2	
8	0.6050	1502	192.1	68.1.1 1	192.168.1.2	
9	0.6150	78	192 1	68.1.3 1	192.168.1.2	
0	0.0100	10	102.	00.1.0	02.100.1.2	
10	0.6200	<u>58</u>			192.168.1.2	
530.0				68.1.1 1	192.168.1.2	I traffic Z _s (t) Itation metho
1000			192.1	68.1.1 1	192.168.1.2 Tranformed by de-fragmen	
5302		<u>58</u>	192.1	68.1.1 1	192.168.1.2 Tranformed by de-fragmen	Destination
5307		58 No.	192.1	68.1.1 1	192.168.1.2 Tranformed by de-fragmen size Source	Destination 192.168.1.2
5302		58 No. 1	192.1 Time 0.0000	68.1.1 1 Packet : 70	192.168.1.2 Tranformed by de-fragmen size Source 192.168.1.1	Destination 192.168.1.2 192.168.1.2
5302		58 No. 1 2	192.1 Time 0.0000 0.1000	Packet : 70 5046	192.168.1.2 Tranformed by de-fragmen size Source 192.168.1.1 192.168.1.1	Destination 192.168.1.2 192.168.1.2 192.168.1.2 192.168.1.2
5307		58 No. 1 2 6	192.1 Time 0.0000 0.1000 0.3000	68.1.1 1 Packet : 70 <u>5046</u> 65	192.168.1.2 Tranformed by de-fragmen size Source 192.168.1.1 192.168.1.1 192.168.1.1	Destination 192.168.1.2 192.168.1.2 192.168.1.2 192.168.1.2 192.168.1.2

Figure 6: Simple example of captured traffic transformed by the defragmentation method

These results are sorted in the histogram of $X_f(t)$. Similarly, as at the first algorithm, a histogram is obtained for $Y_f(t)$ and the identified pdf, as well as estimating their parameters using EasyFIT [23] tool (Fig. 7).

IV. ALGORITHM VALIDATION

Validation of the developed method was performed using an OPNET Modeler simulation tool [14, 16, 17, 19, 20, 21]. An IP station was used for a simulation model (Fig. 1) with Pareto distribution ($\alpha = 1, k = 26$) as the file size process and Weibull distribution ($\alpha = 0.5, k =$ 0.02) as the file inter-arrival time process for the file generation processes. This simulated network traffic represents the referenced traffic, on which the distribution parameters were estimated. With some modifications to the OPNET process model of the IP station (to perform packet logging functionality), the simulated packets' data of the modeled network traffic was captured, during the simulation run. The captured packet's data includes information about packet time_stamps, and the sources and destinations of IP addresses. These data were sufficient for transformation using the defragmentation method. We calculate histograms and estimate distribution parameters on the transformed captured traffic. Any discrepancies between histograms and the chosen distribution are evaluated, using different goodness of fit tests, such as Kolmogorov-Smirnov, Anderson-Darling and Chi-Square [11-13].

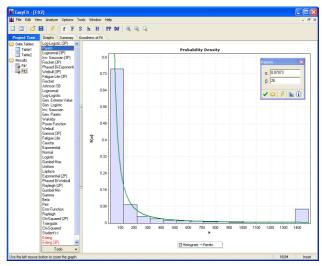


Figure 7: Histogram of a packet size process regarding captured packets, and suitable distribution with estimated parameters in EasyFit fitting tool.

V. CONCLUSION

This paper presents a comparison between two developed algorithms for the calculation of histograms regarding the processes which generate files, presenting a source of packet traffic. Continuous pdf's parameters, using some of the fitting tools, are later identified from these histograms. These parameters can, in-turn, be used as the parameters of a traffic generator in a simulation program.

The results of both algorithms have negligible statistical differences, so on this level the primacy has given way to a second algorithm, due to its simplicity and far shorter execution time. In those cases, where it is desirable to identify those applications that contribute to the packets' traffic, then the first algorithm is easier to expand, in such way that enables simultaneous statistical calculation of any application involved in packet traffic. The second algorithm has no such ability. It can calculate a statistic of any application only if a measurement tool, i.e. sniffer, from all the traffic filtered out only traffic originated in the selected application.

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